

K-Area: An Efficient Approach to Approximate the Spatial Boundaries of Mobility Data with k-Anonymity

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About me - Maël Gassmann



- Swiss
- French native speaker; German
- BSc in Computer Science in 2022 at the BFH
- MSc in Engineering 3rd semester at the BFH
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Summary

1. Introduction
2. State of the Art
3. Concept
4. Conclusion

Introduction

- Mobility datasets
 - Complex
 - Potent in utility
- Anonymity is a **must**
- Privacy preserving methods:
 - Computationally expensive
 - Use case specific

Objectives:

- Fast assessment
- Not use case specific
- Reduce the problem space for privacy enhancing methods

State of the Art

Mitigation:

- Based on Heuristics
- No theoretical or provable guarantees
- Examples: *swapping*, *obfuscation*, *spacial cloaking* or *segmentation*

Indistinguishability:

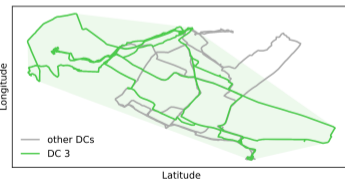
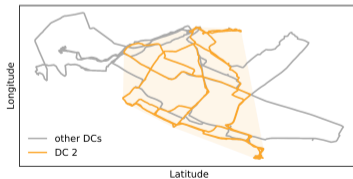
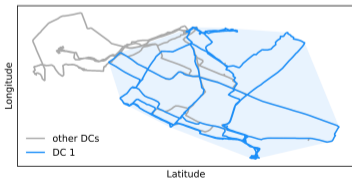
- Anonymity measured in terms of how distinguishable are each data collectors
- Filtering out singularities → better anonymity
- Examples: *k-anonymity*, ***k-area***

Uninformativeness:

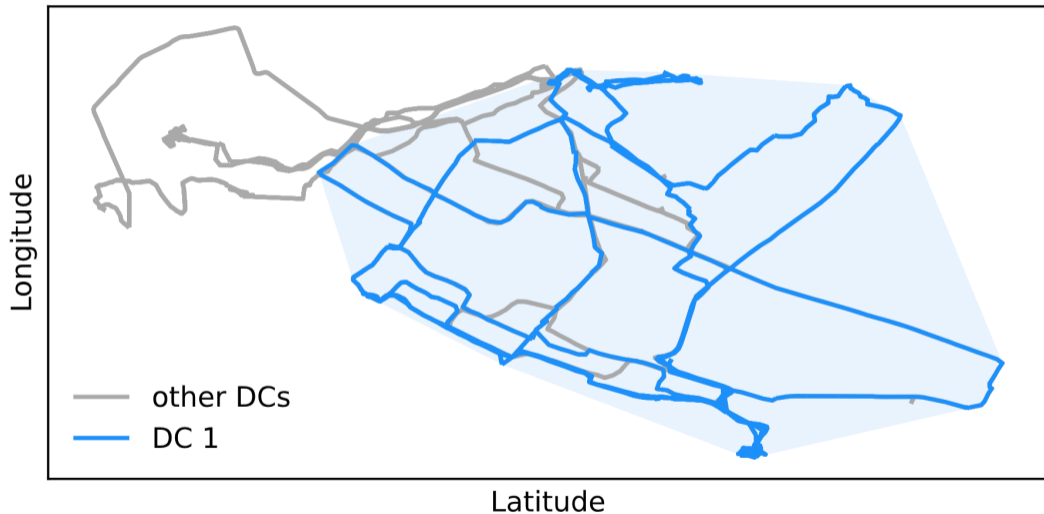
- Anonymity measured in terms of how much information is held per data collectors
- The less of difference → the better the anonymity
- Example: *differential privacy*

Concept

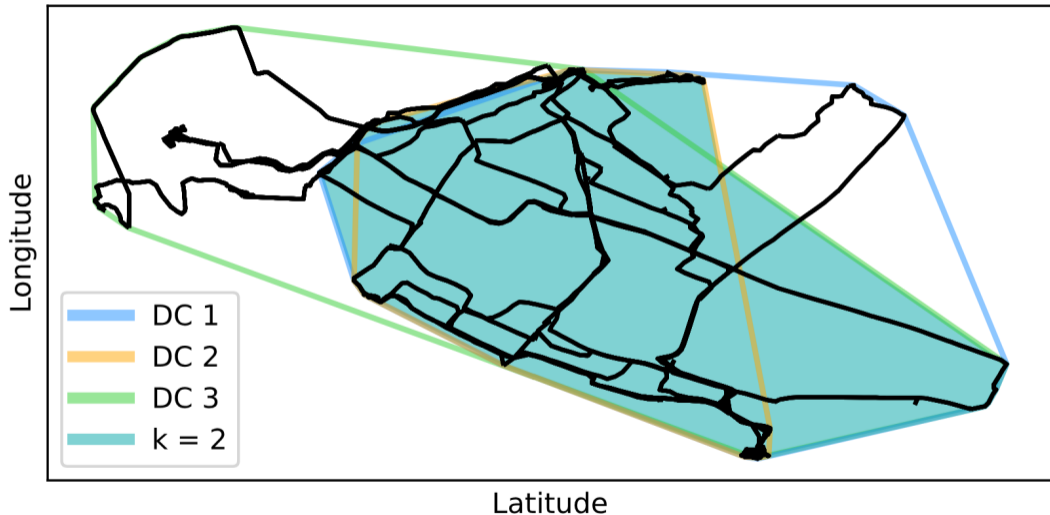
- Mobility dataset as a finite set of:
 - ▣ GPS points
 - ▣ Data collector identifiers
- Always has clear spacio-temporal bounds
- Data Collector spacio-temporal bounds
- K-Area is based on this and strives to reduce them further



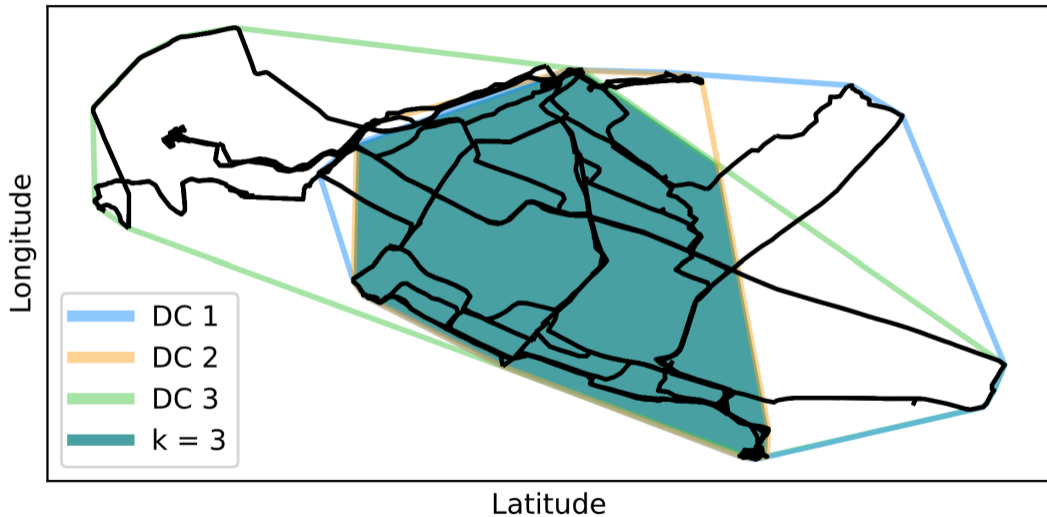
Concept - Polygon



Concept - K-Area



Concept - K-Area



Concept - Algorithm

$$A_k = \bigcup (p_{i_1} \cap p_{i_2} \cap \dots \cap p_{i_k})$$

for all $p_{i_1}, p_{i_2}, \dots, p_{i_k} \in P$
with $i_1 \neq i_2 \neq \dots \neq i_k$ and $k \geq 2$

Concept - Algorithm

```
1: function K-Area( $P, k$ )
2:    $A \leftarrow$  array of size  $k$  with each entry as an empty set of polygons
3:   for all  $p_u \in P$  do
4:     for  $i \leftarrow k$  to 1 do
5:       if  $i > 1$  then
6:          $A[i] \leftarrow A[i] \cup (p_u \cap A[i - 1])$ 
7:       else
8:          $A[i] \leftarrow A[i] \cup p_u$  ▷  $i = 1$ 
9:       end if
10:    end for
11:  end for
12:  return  $A$ 
13: end function
```

Concept - Use cases

- Set a k-anonymity condition and run the algorithm periodically while collecting data. Allowing for better management of surveys during their production.
- Generate heat-maps to visualize the readiness of a dataset and where data could be lacking.
- To be used as a pre-processing step before running computational expensive algorithms.

Conclusion

- A solution was proposed:
 - To quickly assess a dataset
 - That can be used to reduce the problem space of privacy enhancing methods
 - Is not use case specific
 - Discards superfluous parts
- Based on geometric operations
- Polynomial order

Future work:

- Different shapes for the spacial bounds
- Researches concerning its application in 3 dimensions

Questions

Questions ?